# Virtual Network Subsystem

### **Use-Cases**

- Creating virtual SDN networks for tenants
- Slicing regions of networks for use by different tenants (M-CORD)
- Federation exposing abstracted view to peer/parent controllers (E-CORD)

# General Approach

There are many different ways through which network virtualization can be offered to ONOS applications. On one end of the spectrum is a scheme used by OpenStack where virtual network is exposed as a topology-less set of end-station hosts with have inherent L2/L3 connectivity provisioned between them; this is usually accomplished via overlay such as VLAN, VxLAN or GRE tunnels. On the other end of the spectrum is a scheme employed by OpenVirtex (OVX) where virtual networks are exposed as SDN networks without any inherent connectivity; in order to program connectivity, one must use OpenFlow to control and program virtualized network devices.

As part of OpenStack and OPNFV integration, ONOS already provides means to provision OpenStack-style virtual networks via overlays. The Virtual Tenant Network (VTN) application offers these capabilities.

In order for ONOS to offer applications and users the ability to program their own virtual networks, ONOS needs to also support the ability to create such SDN-capable virtual networks. This is where the virtual network subsystem comes in along with its set of first-class virtual network model constructs, e.g. virtual device, virtual link, virtual port.

# **Basic Concepts**

Virtual network belongs to a tenant

In order to allow tracking virtual networks to tenants, ONOS supports a notion of a tenant. However, this notion is purposefully kept very shallow and a *tenant is represented merely by a tenant ID*. This way, the tenant ID effectively acts as a foreign key that can be used to tie information with other external systems, like OpenStack, which have been designed to keep deeper information about tenants.

#### Virtual network has virtual devices and links

In order to represent topology, ONOS *virtual networks form a graph consisting of virtual devices and virtual links*, derivatives of devices and links. In this manner, the virtual network topology graph can be traversed and treated in the same manner as a physical network topology graph.

#### Virtual device has virtual ports

Virtual devices consist of virtual ports, each of which is *realized by a port from the underlying network*. Therefore, the only ties between the virtual network and the underlying physical one are the virtual port to physical port mappings; *a virtual device is effectively a slicing/splicing of devices from the underlying network*.

TODO: Add a diagram showing possible scenarios with virtual to physical mappings, including interior ports as well as edge ports and virtual hosts.

Ali: Virtual hosts should continue to belong to the virtual network regardless of their location.

#### Virtual devices & links are elastic entities

Since the definition of virtual devices and virtual links are expressed merely via their constituent port entities, they are not backed by a specific path and therefore do not necessarily depend on a specific structure of the underlying network.

Q: Should a virtual port be represented by a group of ports? Group of infra ports? Ali: Suggestion to leave this for later.

Ali: Allow ability to specify "mode" when creating links, e.g. fail-safe, specific paths, etc. The idea is to allow different recovery models.

# **Basic Operation**

Since virtual networks have topology, apps can traverse topology graph using a regular TopologyService and/or PathService capabilities, though only in the context narrowed to a specific virtual network.

Due to the elastic nature of the binding between the virtual network and the underlying network, the virtual network topology changes only in response to SCC changes in the underlying network, meaning when the continuity of the underlying network changes. Therefore, the virtual devices and links go up/down only when potential for connectivity and availability changes in the underlay.

Similarly, because virtual networks have no inherent connectivity, no action needs to be taken when a network is created. Instead, the act of programming the virtual network drives action on

the underlying network; any logical tunnels need to be created on the underlying network only when connectivity is requested at the virtual network topology.

#### Virtual Network Admin Service

To permit construction of virtual network, a privileged VirtualNetworkAdminService is provided offering the following operations:

- Register/unregister tenants (TenantId)
- Create/remove networks (for tenant)
- Create/remove device/link/host (for network)
- Create/remove port (for device)
  - This also binds the port to the physical network

#### Virtual Network Service

To query the virtual network inventory, the following list of operations is provided by the VirtualNetworkService interface:

- List of tenants (TenantId)
- List of networks (by tenant)
- List of devices/links/hosts (by network)
- List of ports (by device)

Apps can also obtain "narrowed" services using which they can operate on the virtual network just as they would on a regular physical network. These services are confined to the only those entities comprising the given virtual network and include the following:

- DeviceService, LinkService, HostService, TopologyService, PathService
- IntentService

### Virtual Network Provider & Service

In order to allow the virtualization mechanism to vary, the virtual network subsystem purposefully separates the modeling and tracking of the virtual networks apart from the mechanism used to actually realize the networks on the data-plane. This separation is enforced using the VirtualNetworkProvider and VirtualNetworkProviderService interfaces. These interfaces are still in flux as we experiment with different mechanisms in order to find the minimal necessary interface with sufficient signaling between the virtual network subsystem and data-plane providers.

In general, the responsibilities of a virtual network provider include the following:

- Provide elastic connectivity "tunnels"
  - for link traversal
  - o for device "backplane" traversal, i.e. port connectivity
  - set up "on-demand" in response to programming edicts
- Maintain tunnels in face of failures

- device and link failures in underlying network
- congestion in underlying network
- Report events to provider service
  - irrecoverable network outages

### Virtual Networks & Apps

#### On-Platform Apps

Because ONOS on-platform apps are effectively core-extensions that are deployed in symmetric fashion across the entire ONOS cluster and because the base OSGi platform is capable of only limited means of compute/memory resource isolation, ONOS applications cannot be deployed on behalf of specific tenants or specific virtual networks. As such they cannot be tenant or network-specific.

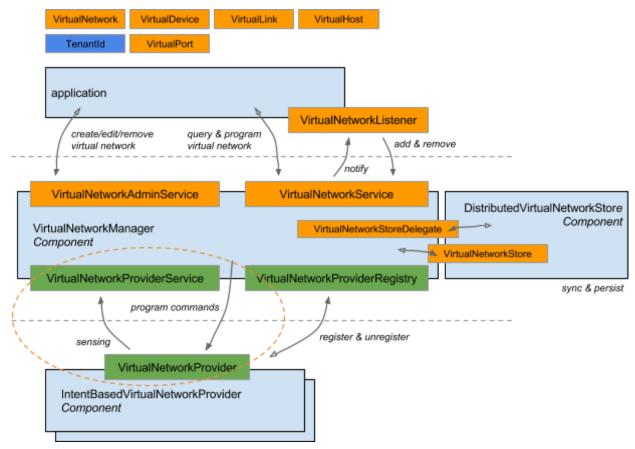
On-platform apps can, however, be tenant and network-aware, meaning that they can operate on behalf of all tenants or networks in their respective contexts.

#### Off-Platform Apps

Off-platform applications, which by definition run external to the ONOS OSGi platform, and which do not have to adhere to the symmetric deployment constraint are free to be tenant-specific or virtual network-specific and operate using the ONOS REST API to create and control virtual networks. Because such apps run on a separate platform, compute & memory resource isolation is outside of ONOS' scope and would typically be done using OS specific means or by using dedicated hardware.

It should be noted that use of REST API for tenants to interact with their networks is predicated on existence of a scoped authorization scheme, which is one of the features on the roadmap - see section below.

# High-Level Design



Note: The circled area continues to be subject of design and experimentation.

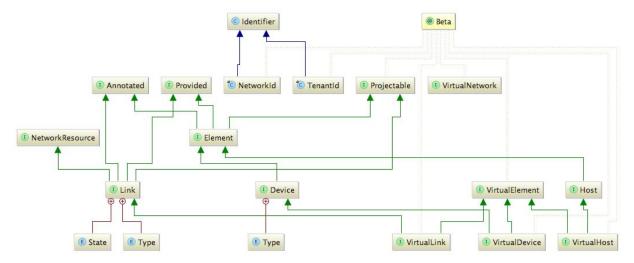
The general idea of this design is to provide a clear separation between modeling of the virtual networks and their realization on the physical underlays. The model and the methods to manipulate the structure of the virtual network remains the same regardless of how the virtual networks are realized.

In order to provide this separation, the VirtualNetworkProvider and VirtualNetworkProviderService interfaces, which provide the basis for this boundary, need to be design with the minimum necessary surface to allow these parts of the system to communicate, but without introducing dependencies on specific tunneling mechanisms that may be used to implement the virtual network providers. This continues to be an area of active development.

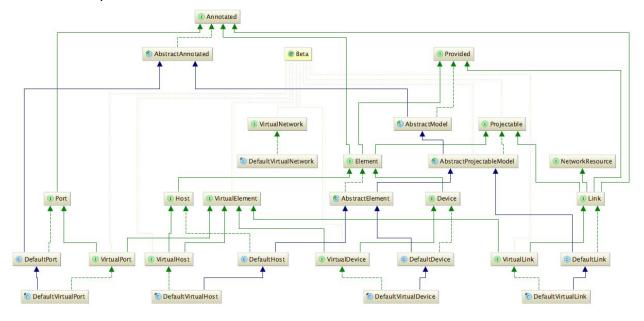
The code is located in the incubator portion of the source tree and under the org.onosproject.incubator.net.virtual Java package hierarchy. To reflect the fact that the APIs are still in flux, the model and service interface code is marked with @Beta annotation.

# Model Object Class Hierarchy

The following UML diagram depicts the interface hierarchy of the various virtual network model entities and their relationships with the rest of the ONOS core model.



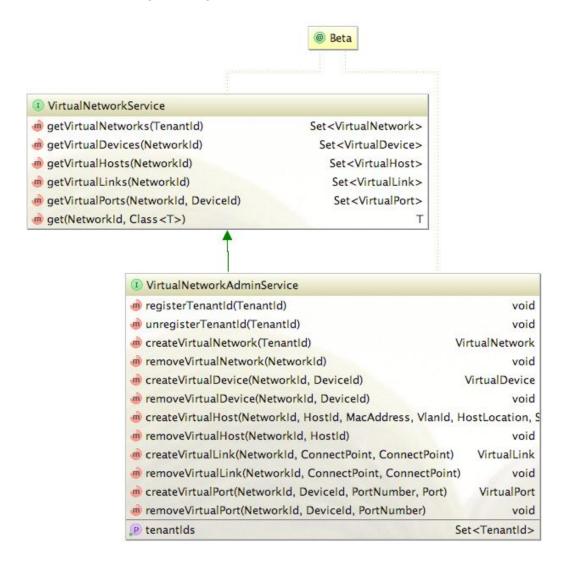
Note that the virtual elements such as virtual device, link and host descend from their core model counterparts.



### Virtual Network Service & Admin Service

The following depicts the functionality available from the VirtualNetworkService and the administrative extension VirtualNetworkAdminService. Note that the create/edit/delete

functionality is available solely through the administrative interface, whereas the query capabilities are available through the regular service interface.



In order to operate on the virtual network, applications are expected to use the T VirtualNetworkService.get(NetworkId, Class<T>) method to obtain a reference to the specified class of service, which has been narrowed to operate solely in the confines of the specified virtual network. In this way, an application can use these services to operate on the virtual network in exactly the same way as it would if it were operating on a physical network.

## Virtual Gateways

An expected requirement is to allow customers to connect virtual networks to internet and to other virtual networks via a gateway.

There are several ways in which this capability can be accomplished, but presently no specific approach has been selected yet.

### **REST API**

Presently, the virtual network subsystem exposes a REST API end-point to allow off-platform applications to query, create, edit and remove virtual networks and to program them via intents.

Presently, the end-point does not provide any authorization capabilities to limit tenants to operate only on their own networks, but this functionality is expected to be added in the future as the implementation of the subsystem matures.

### CLI

A set of CLI commands have been provided to allow the administrator to use the ONOS shell to query, create, edit and remove virtual networks. The commands are listed below:

- vnet-tenants list tenant ids
- vnet-add-tenant add a new tenant
- vnet-remove-tenant remove a tenant
- vnets list virtual networks
- vnet-create networkld tenantld create a new virtual network
- vnet-remove networkld remove virtual network
- vnet-devices list all virtual devices in a virtual network
- vnet-create-device networkId deviceId create a virtual device in a virtual network
- vnet-remove-device networkId deviceId remove a virtual device from a virtual network
- vnet-ports list all virtual ports in a virtual network
- vnet-create-port networkld deviceld/portNumber physDevice/physPortNumber create a virtual port
- vnet-remove-port networkld deviceld/portNumber
- vnet-links list all virtual links in a virtual network
- vnet-create-link networkld srcDeviceld/portNumber dstDeviceld/portNumber [bidirectional]
- vnet-remove-link networkld srcDeviceId/portNumber dstDeviceId/portNumber [bidirectional]

### **Limitations & Caveats**

### OpenFlow programming

Current approach does not expose a device-specific configuration of programming API for virtual devices. However, the design does not preclude it and the long-term goal is to enable virtual device programming via Flow Objectives and based on that expose OpenFlow interface through which virtual devices can be controlled via network controllers that support OpenFlow protocol.

### Virtual Network Gateways

Design for connecting virtual networks to each other and to internet has not been selected/specified yet.

#### **REST API Authorization**

Current REST API has no authorization mechanism.

# Summary

- Virtual networks have no inherent connectivity
- Instead, connectivity is programmed via intents (in future flow objectives)
- Connectivity for link & device traversal are setup only when connectivity is programmed
- Virtual port is realized by underlying port
- Virtual networks must be able to share ports
  - to share exterior ports, ingress traffic must be tagged
  - o to share interior ports, pass-through traffic is tunneled

# Pending Questions, Suggestions & TODOs

Thomas: Add diagrams showing various forms of mappings between virtual & physical.

Aihua: Are the virtual networks L2/L3 or are they capable of virtualizing L0/L1 as well, e.g. transport network virtualization?

Bill: Advertise the document & video to solicit feedback and comments.

Bill: Can we get use cases from {E,R,M}-CORD how this will be used?

E.g. Running rcord and ecord together; MVNO; how to get the same functionality as VTN

Yuta: E-CORD has a crude big switch app that is used to expose DC to parent cluster. Maybe we should move to the virtualization abstractions.

Ali: this is more of a federation controller, OF isn't desired here, this is more of an API call interaction